

POSCAT Seminar 3-1 : Brute Force

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Topic

- Topic today
 - Brute Force
 - Time Complexity
 - Algorithm Design
 - Implementation



Brute Force

- 모든 경우를 고려하여 답을 찾자
 - 모든 경우를 고려하는 알고리즘을 생각하고 // **Verification**
 - 그 경우의 수가 실제로 얼마나 되는지를 계산 // **Efficiency**
- 가장 기본적인 접근
 - 어떤 문제를 받던 **가장 먼저** Brute Force하게 접근
 - 경우의 수가 **충분히 적다면** 단순히 Brute Force하게 풀면 된다
 - 단순하게 생각 !



Brute Force

- Example

N개의 숫자가 주어졌을 때, 연속부분 최대합을 구하라

2	-4	1	2	-1	3	-7	2	1	-4	-2
---	----	---	---	----	---	----	---	---	----	----



Brute Force

- Example

N개의 숫자가 주어졌을 때, 연속부분 최대합을 구하라

2	-4	1	2	-1	3	-7	2	1	-4	-2
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Brute Force

- Example

가장 먼저 Brute Force하게 생각을 해 보자

2	-4	1	2	-1	3	-7	2	1	-4	-2
---	----	---	---	----	---	----	---	---	----	----



Brute Force

- Example

가장 먼저 Brute Force하게 생각을 해 보자
→ 모든 경우를 고려한다는 뜻은 무엇인가?

2	-4	1	2	-1	3	-7	2	1	-4	-2
---	----	---	---	----	---	----	---	---	----	----



Brute Force

■ Example

가장 먼저 Brute Force하게 생각을 해 보자

→ 모든 경우를 고려한다는 뜻은 무엇인가?

→ 모든 연속부분의 합을 구해보고 그 중에 최댓값 선택 !

2	-4	1	2	-1	3	-7	2	1	-4	-2
---	----	---	---	----	---	----	---	---	----	----



Brute Force

- Example

모든 연속부분의 합은 어떻게 구할 것인가 ?

2	-4	1	2	-1	3	-7	2	1	-4	-2
---	----	---	---	----	---	----	---	---	----	----



Brute Force

- Example

모든 연속부분의 합은 어떻게 구할 것인가 ?

연속 부분의 시작점과 끝점을 정하면 구간이 정해짐

2	-4	1	2	-1	3	-7	2	1	-4	-2
---	----	---	---	----	---	----	---	---	----	----



Brute Force

- Example

모든 연속부분의 합은 어떻게 구할 것인가 ?

연속 부분의 시작점과 끝점을 정하면 구간이 정해짐

가능한 모든 시작점과 끝점을 선택하여 합을 구하면 되는군

2	-4	1	2	-1	3	-7	2	1	-4	-2
---	----	---	---	----	---	----	---	---	----	----



Brute Force

- Example

모든 연속부분의 합은 어떻게 구할 것인가 ?

연속 부분의 시작점과 끝점을 정하면 구간이 정해짐

가능한 모든 시작점과 끝점을 선택하여 합을 구하면 되는군

2	-4	1	2	-1	3	-7	2	1	-4	-2
---	----	---	---	----	---	----	---	---	----	----

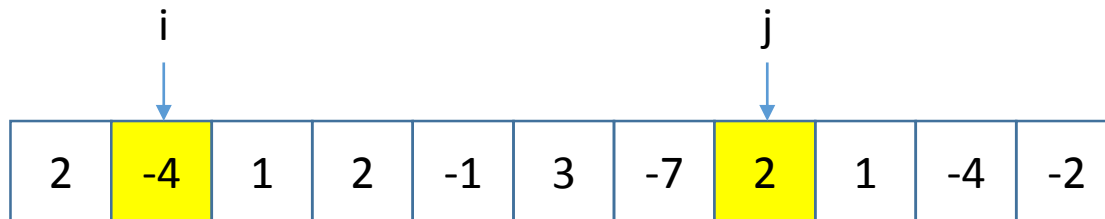
More concretely !



Brute Force

- Example

i를 시작점이라 하고, j를 끝점이라 하자

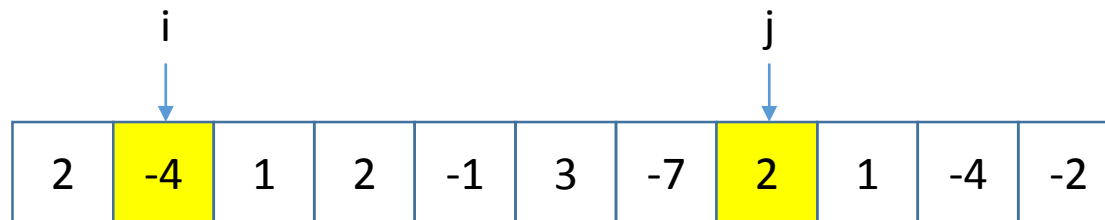


Brute Force

- Example

i 를 시작점이라 하고, j 를 끝점이라 하자

모든 가능한 i 와 j 에 대하여 $D[i] \sim D[j]$ 의 합을 구하면 됨

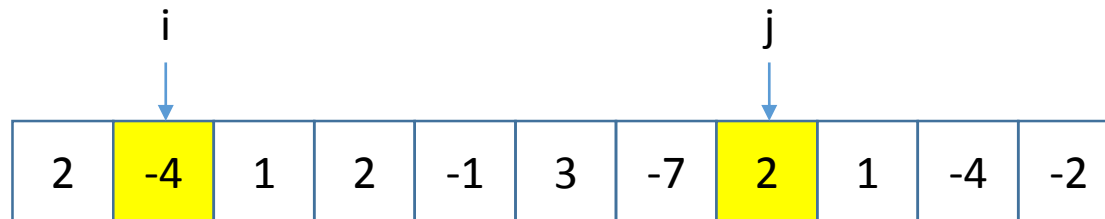


Brute Force

- Example

i를 시작점이라 하고, j를 끝점이라 하자

모든 가능한 i와 j에 대하여 $D[i] \sim D[j]$ 의 합을 구하면 됨



GOOD !



Brute Force

```
1 int result = -inf; // -infinity
2 for(int i=0;i<n;i++){
3     for(int j=i;j<n;j++){
4         int sum = 0;
5         for(int k=i;k<=j;k++) sum += D[k];
6         result = max(result, sum);
7     }
8 }
```

Verification : Is this algorithm **true** ?

Efficiency : How long does it **take** ?



Brute Force

```
1 int result = -inf; // -infinity
2 for(int i=0;i<n;i++){
3     for(int j=i;j<n;j++){
4         int sum = 0;
5         for(int k=i;k<=j;k++) sum += D[k];
6         result = max(result, sum);
7     }
8 }
```

Verification : Is this algorithm **true** ? We consider all the cases

Efficiency : How long does it **take** ? $O(n^3)$



Brute Force

```
1 int result = -inf; // -infinity
2 for(int i=0;i<n;i++){
3     for(int j=i;j<n;j++){
4         int sum = 0;
5         for(int k=i;k<=j;k++) sum += D[k];
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```

Verification : Is this algorithm **true** ? We consider all the cases

Efficiency : How long does it **take** ? $O(n^3)$

Is this algorithm efficient to solve this problem ?



Rough Criteria

- 약 1억 번의 연산이 1초라고 생각
 - Time Complexity 를 계산했을 때 '1억'이면 대략 1초
 - O-notation 자체가 **근사**이기 때문에 **완전히 믿기는 힘들**
 - $O(n^3)$ VS $O(99999999n^2)$
- 의심되면 좀 더 정확하게 계산
 - Constant factor 까지 좀 더 **세밀하게** 고려하여 계산
 - 어느 정도의 **감**이 있는 것이 굉장히 중요
- 보통 의도한 솔루션으로 풀면 ACCEPT
 - 계산 결과가 아슬아슬하면 대부분의 경우 **방법이 틀림**
 - 알고 보니 **해법 자체**가 아슬아슬한 문제라면 **굉장히 어려운 문제!!**

out of scope ☺



Brute Force

```
1 int result = -inf; // -infinity
2 for(int i=0;i<n;i++){
3     for(int j=i;j<n;j++){
4         int sum = 0;
5         for(int k=i;k<=j;k++) sum += D[k];
6         result = max(result, sum);
7     }
8 }
```

Verification : Is this algorithm **true** ? We consider all the cases

Efficiency : How long does it **take** ? $O(n^3)$

Is this algorithm efficient to solve this problem ?
Suppose $n \leq 100$ and Time Limit = 1s



Brute Force

```
1 int result = -inf; // -infinity
2 for(int i=0;i<n;i++){
3     for(int j=i;j<n;j++){
4         int sum = 0;
5         for(int k=i;k<=j;k++) sum += D[k];
6         result = max(result, sum);
7     }
8 }
```

Verification : Is this algorithm **true** ? We consider all the cases

Efficiency : How long does it **take** ? $O(n^3)$

Is this algorithm efficient to solve this problem ?

Suppose $n \leq 100$ and Time Limit = 1s

$$n^3 = 1,000,000 \leq 100,000,000$$



Brute Force

```
1 int result = -inf; // -infinity
2 for(int i=0;i<n;i++){
3     for(int j=i;j<n;j++){
4         int sum = 0;
5         for(int k=i;k<=j;k++) sum += D[k];
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```

Verification : Is this algorithm **true** ? We consider all the cases

Efficiency : How long does it **take** ? $O(n^3)$

Is this algorithm efficient to solve this problem ?

Suppose $n \leq 100$ and Time Limit = 1s

$$n^3 = 1,000,000 \leq 100,000,000$$

GOOD!!



Brute Force

```
1 int result = -inf; // -infinity
2 for(int i=0;i<n;i++){
3     for(int j=i;j<n;j++){
4         int sum = 0;
5         for(int k=i;k<=j;k++) sum += D[k];
6         result = max(result, sum);
7     }
8 }
```

Verification : Is this algorithm **true** ? We consider all the cases

Efficiency : How long does it **take** ? $O(n^3)$

Is this algorithm efficient to solve this problem ?
Suppose $n \leq 1000$ and Time Limit = 1s



Brute Force

```
1 int result = -inf; // -infinity
2 for(int i=0;i<n;i++){
3     for(int j=i;j<n;j++){
4         int sum = 0;
5         for(int k=i;k<=j;k++) sum += D[k];
6         result = max(result, sum);
7     }
8 }
```

Verification : Is this algorithm **true** ? We consider all the cases

Efficiency : How long does it **take** ? $O(n^3)$

Is this algorithm efficient to solve this problem ?

Suppose $n \leq 1000$ and Time Limit = 1s

$$n^3 = 1,000,000,000 > 100,000,000$$



Brute Force

```
1 int result = -inf; // -infinity
2 for(int i=0;i<n;i++){
3     for(int j=i;j<n;j++){
4         int sum = 0;
5         for(int k=i;k<=j;k++) sum += D[k];
6         result = max(result, sum);
7     }
8 }
```

Verification : Is this algorithm **true** ? We consider all the cases

Efficiency : How long does it **take** ? $O(n^3)$

Is this algorithm efficient to solve this problem ?

Suppose $n \leq 1000$ and Time Limit = 1s

$$n^3 = 1,000,000,000 > 100,000,000$$

BAD, this solution is wrong



Brute Force

```
1 int result = -inf; // -infinity
2 for(int i=0;i<n;i++){
3     for(int j=i;j<n;j++){
4         int sum = 0;
5         for(int k=i;k<=j;k++) sum += D[k];
6         result = max(result, sum);
7     }
8 }
```

Verification : Is this algorithm **true** ? We consider all the cases

Efficiency : How long does it **take** ? $O(n^3)$

Is this algorithm efficient to solve this problem ?
Suppose $n \leq 470$ and Time Limit = 1s



Brute Force

```
1 int result = -inf; // -infinity
2 for(int i=0;i<n;i++){
3     for(int j=i;j<n;j++){
4         int sum = 0;
5         for(int k=i;k<=j;k++) sum += D[k];
6         result = max(result, sum);
7     }
8 }
```

Verification : Is this algorithm **true** ? We consider all the cases

Efficiency : How long does it **take** ? $O(n^3)$

Is this algorithm efficient to solve this problem ?

Suppose $n \leq 470$ and Time Limit = 1s

$$n^3 = 103,823,000 > 100,000,000$$



Brute Force

```
1 int result = -inf; // -infinity
2 for(int i=0;i<n;i++){
3     for(int j=i;j<n;j++){
4         int sum = 0;
5         for(int k=i;k<=j;k++) sum += D[k];
6         result = max(result, sum);
7     }
8 }
```

Verification : Is this algorithm **true** ? We consider all the cases

Efficiency : How long does it **take** ? $O(n^3)$

BAD?

Is this algorithm efficient to solve this problem ?

Suppose $n \leq 470$ and Time Limit = 1s

$$n^3 = 103,823,000 > 100,000,000$$



Brute Force

```
1 int result = -inf; // -infinity
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6         result = max(result, sum);
7     }
8 }
```

Verification : Is this algorithm **true** ? We consider all the cases

Efficiency : How long does it **take** ? $O(n^3)$

BAD?

No. n^3 is over-estimated

In fact, It takes less than $\frac{n^3}{2}$

Is this algorithm efficient to solve this problem ?

Suppose $n \leq 470$ and Time Limit = 1s

$$n^3 = 103,823,000 > 100,000,000$$



Brute Force

```
1 int result = -inf; // -infinity
2 for(int i=0;i<n;i++){
3     for(int j=i;j<n;j++){
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```

Verification : Is this algorithm **true** ? We consider all the cases

Efficiency : How long does it **take** ? $O(n^3)$

BAD?

No. n^3 is over-estimated

In fact, It takes less than $\frac{n^3}{2}$

Is this algorithm efficient to solve this problem ?

Suppose $n \leq 470$ and Time Limit = 1s

$$\frac{n^3}{2} = 51,911,500 \leq 100,000,000$$



Brute Force

```
1 int result = -inf; // -infinity
2 for(int i=0;i<n;i++){
3     for(int j=i;j<n;j++){
4         int sum = 0;
5         for(int k=i;k<=j;k++) sum += D[k];
6         result = max(result, sum);
7     }
8 }
```

Verification : Is this algorithm **true** ? We consider all the cases

Efficiency : How long does it **take** ? $O(n^3)$

GOOD ENOUGH!

BAD?

No. n^3 is over-estimated

In fact, It takes less than $\frac{n^3}{2}$

Is this algorithm efficient to solve this problem ?

Suppose $n \leq 470$ and Time Limit = 1s

$$\frac{n^3}{2} = 51,911,500 \leq 100,000,000$$



It is hard to calculate perfectly

- 완벽하게 계산하는 것은 굉장히 어려움
 - 당장 이전 슬라이드의 3중 for문이 $\frac{n^3}{2}$ 보다는 확실히 덜 걸림
 - 문제를 많이 풀다가 보면 감이 생깁니다
- 아슬아슬하면 다른 솔루션을 찾자
 - 아까도 이야기 했지만, 아슬아슬하면 솔루션이 틀린 경우
 - 다른 방법이 없을지를 생각하는 것이 좋다

out of scope ☺



Brute Force

```
1 int result = -inf; // -infinity
2 for(int i=0;i<n;i++){
3     for(int j=i;j<n;j++){
4         int sum = 0;
5         for(int k=i;k<=j;k++) sum += D[k];
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7     }
8 }
```

Verification : Is this algorithm **true** ? We consider all the cases

Efficiency : How long does it **take** ? $O(n^3)$

Is this algorithm efficient to solve this problem ?

Suppose $n \leq 470$ and Time Limit = 1s

$$n^3 = 103,823,000 > 100,000,000$$

BAD?



Brute Force

```
1 int result = -inf; // -infinity
2 for(int i=0;i<n;i++){
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Verification : Is this algorithm **true** ? We consider all the cases

Efficiency : How long does it **take** ? $O(n^3)$

Is this algorithm efficient to solve this problem ?

Suppose $n \leq 470$ and Time Limit = 1s

$$n^3 = 103,823,000 > 100,000,000$$

BAD?

Hmm... Is there another solution ?



Brute Force

We can remove this loop !

```
1 int result = -inf; // -infinity
2 for(int i=0;i<n;i++){
3     for(int j=i;j<n;j++){
4         int sum = 0;
5         for(int k=i;k<=j;k++) sum += D[k];
6         result = max(result, sum);
7     }
8 }
```

Verification : Is this algorithm **true** ? We consider all the cases

Efficiency : How long does it **take** ? $O(n^3)$

Is this algorithm efficient to solve this problem ?

Suppose $n \leq 470$ and Time Limit = 1s

$$n^3 = 103,823,000 > 100,000,000$$

BAD?

Hmm... Is there another solution ?



Brute Force

```
1 S[0] = D[0];
2 for(int i=1;i<n;i++) S[i] = S[i-1] + D[i];
3
4 int result = -inf;
5
6 for(int i=0;i<n;i++){
7     for(int j=i;j<n;j++){
8         int sum = S[j] - S[i-1]; // sum of D[i] ~ D[j]
9         result = max(result, sum);
10    }
11 }
```

Verification : Is this algorithm **true** ?

Efficiency : How long does it **take** ?



Brute Force

```
1 S[0] = D[0];
2 for(int i=1;i<n;i++) S[i] = S[i-1] + D[i];
3
4 int result = -inf;
5
6 for(int i=0;i<n;i++){
7     for(int j=i;j<n;j++){
8         int sum = S[j] - S[i-1]; // sum of D[i] ~ D[j]
9         result = max(result, sum);
10    }
11 }
```

Verification : Is this algorithm **true** ? Still, we consider all the cases

Efficiency : How long does it **take** ? $O(n^2)$



Brute Force

```
1 S[0] = D[0];
2 for(int i=1;i<n;i++) S[i] = S[i-1] + D[i];
3
4 int result = -inf;
5
6 for(int i=0;i<n;i++){
7     for(int j=i;j<n;j++){
8         int sum = S[j] - S[i-1]; // sum of D[i] ~ D[j]
9         result = max(result, sum);
10    }
11 }
```

Verification : Is this algorithm **true** ? Still, we consider all the cases

Efficiency : How long does it **take** ? $O(n^2)$

Is this algorithm efficient to solve this problem ?

Suppose $n \leq 470$ and Time Limit = 1s

$$n^2 = 220,900 \leq 100,000,000$$



Brute Force

```
1 S[0] = D[0];
2 for(int i=1;i<n;i++) S[i] = S[i-1] + D[i];
3
4 int result = -inf;
5
6 for(int i=0;i<n;i++){
7     for(int j=i;j<n;j++){
8         int sum = S[j] - S[i-1]; // sum of D[i] ~ D[j]
9         result = max(result, sum);
10    }
11 }
```

Verification : Is this algorithm **true** ? Still, we consider all the cases

Efficiency : How long does it **take** ? $O(n^2)$

Completely Right !

Is this algorithm efficient to solve this problem ?

Suppose $n \leq 470$ and Time Limit = 1s

$$n^2 = 220,900 \leq 100,000,000$$



Brute Force

- 어떤 문제를 잡건 **시작은 Brute Force !!!!!**
 - **문제를 이해**하는 것을 도와주고
 - **전체 경우의 수**가 얼마나 되는지를 계산하게 하며
 - **솔루션의 모양, 특성**을 찾는 것을 도와준다
- Recursion이 합쳐지면 완벽한 Brute Force
 - **Backtracking** 이 모든 시도를 다 해보는 것
 - Recursion 없이는 모든 경우를 탐색하지 못하는 경우가 있습니다
 - n 개의 숫자 중에서 m 개 뽑기 → We need n for loops
How can you generate n for loops ?



Question ?

- To-do List
 - Basic implementation problems

