



2013 Fall Semester SMP



Day 4 - Recursion

Contents

Data type with I/O

Function

Conditional statement

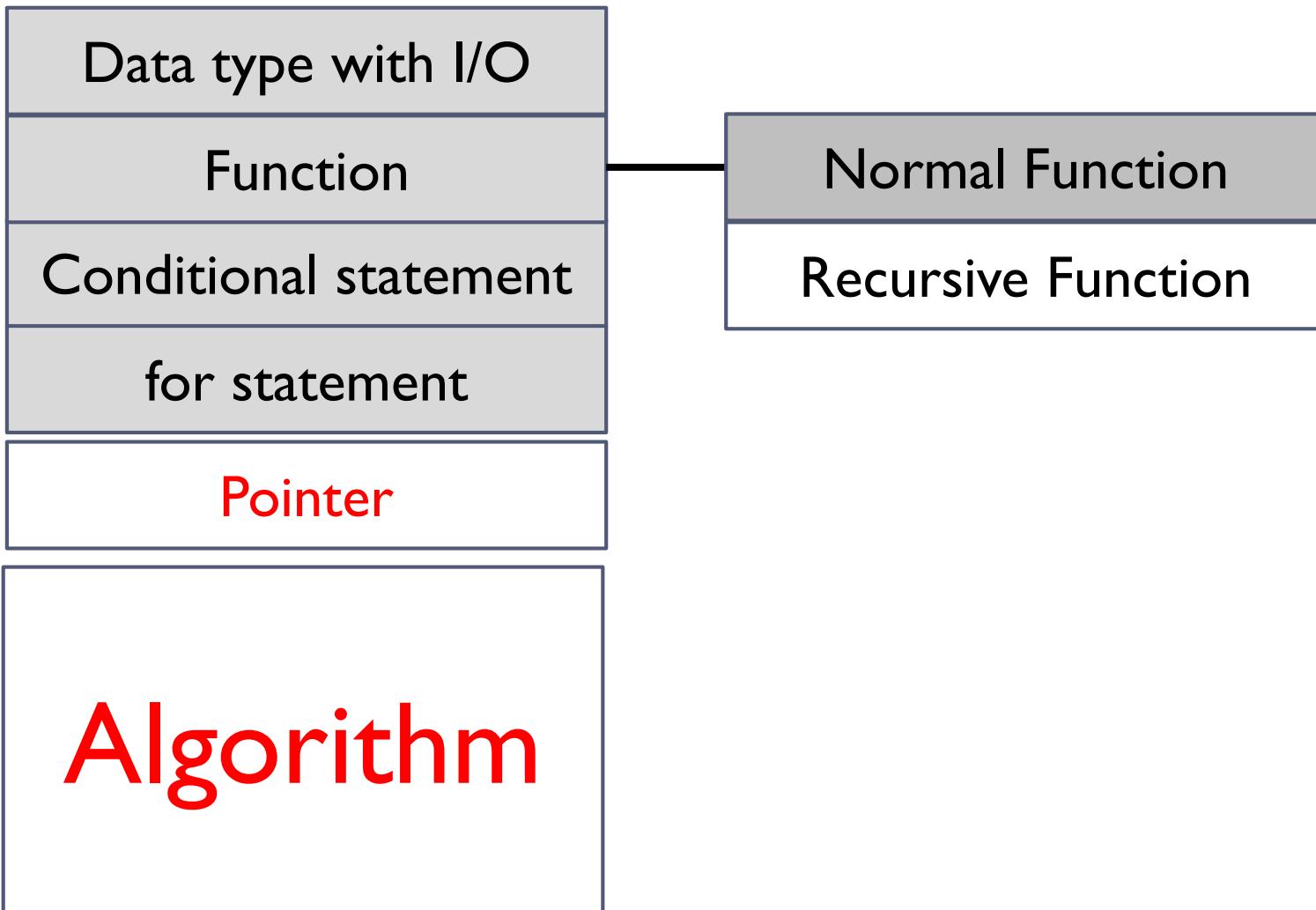
for statement

Pointer

Algorithm



Contents



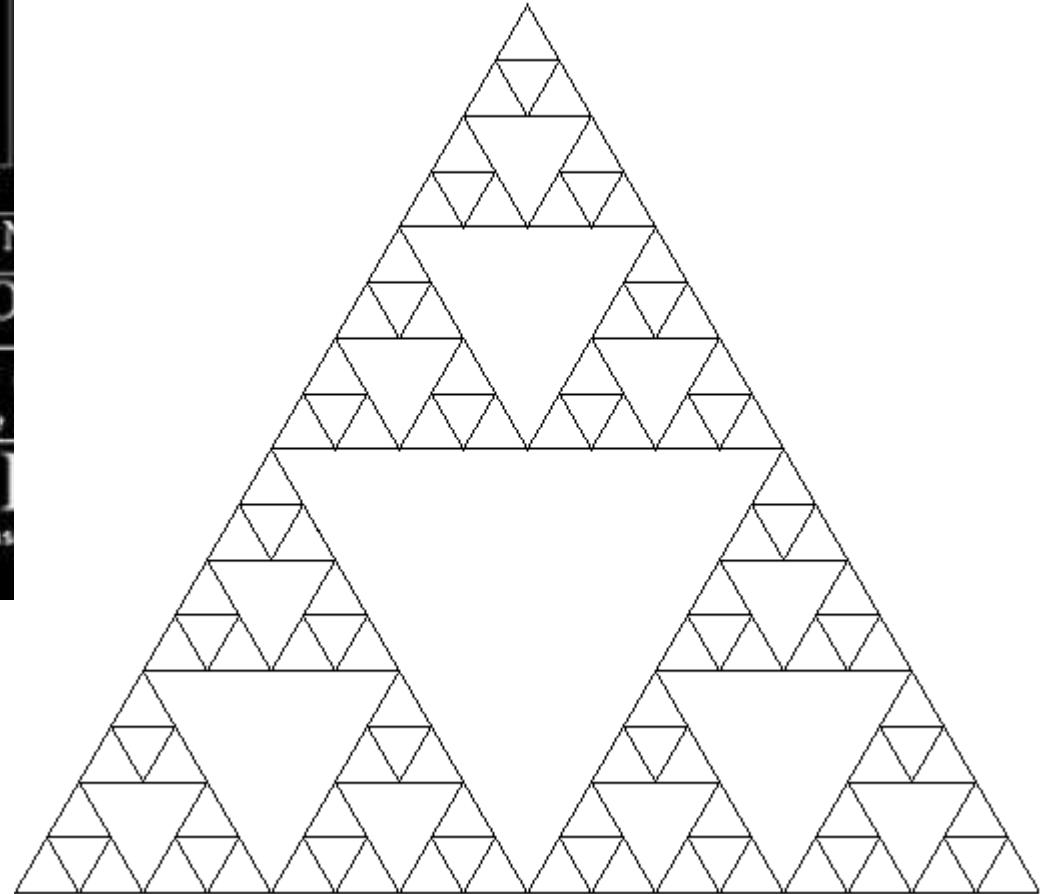
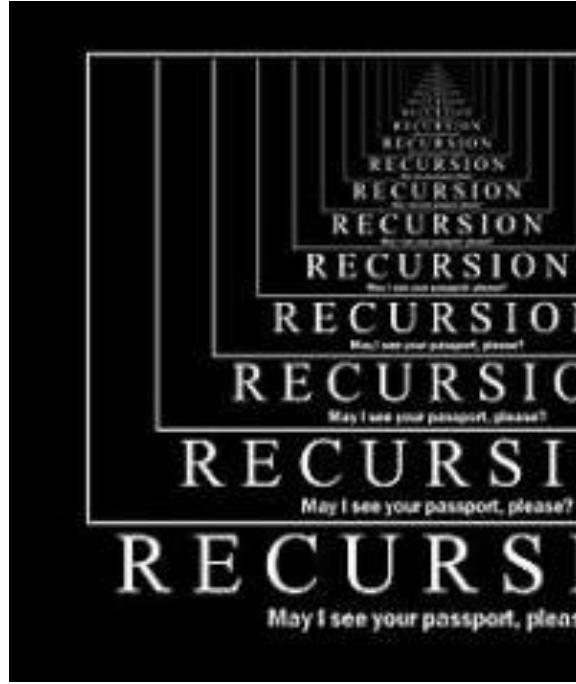
Recursion

- ▶ **Recursion** is the process of repeating items in a **self-similar way** (from wiki)
- ▶ **Recursive function** is implemented in a programming language, whose implementation references itself
 - ▶ 자기 자신을 호출하는 함수
 - ▶ Like..

```
1: #include <cstdio>
2: #include <algorithm>
3: using namespace std;
4: int main(){
5:     main();
6:     return 0;
7: }
```



Recursion



Recursion Example

▶ Fibonacci number

- ▶ 재귀적으로 정의됨 (Recurrence relation)
- ▶ $F_n = F_{n-1} + F_{n-2}$

```
1: #include <cstdio>
2: #include <algorithm>
3: using namespace std;
4: int get_fibo(int x){
5:     if(x == 0) return 1;
6:     else if(x == 1) return 1;
7:     else return get_fibo(x-1) + get_fibo(x-2);
8: }
9: int main(){
10:    printf("%d",get_fibo(5));
11:    return 0;
12: }
```

Recursion Example

▶ Factorial

- ▶ 마찬가지로 점화식이 있음
- ▶ $n! = n \times (n - 1)!$

```
1: #include <cstdio>
2: #include <algorithm>
3: using namespace std;
4: int get_fact(int x){
5:     if(x == 0) return 1;
6:     else return x * get_fact(x-1);
7: }
8: int main(){
9:     printf("%d",get_fact(5));
10:    return 0;
11: }
```

Recursion & Induction

- ▶ 재귀는 수학적 귀납법과 매우 비슷하다
 - ▶ 수학적 귀납법의 경우, 명제를 증명할 때 우선
 - ▶ Base condition을 설정하고,
 - ▶ 재귀적인 관계를 이용해서 증명한다
- ▶ Prove that $\sum_{i=1}^n i = \frac{n(n+1)}{2}$ for $n \in N$
 - ▶ Base condition ($n=1$) : $\sum_{i=1}^1 i = 1 = \frac{1(1+1)}{2}$
 - ▶ Suppose that $\sum_{i=1}^k i = \frac{k(k+1)}{2}$ then
$$\sum_{i=1}^{k+1} i = (k + 1) + \boxed{\sum_{i=1}^k i} = (k + 1) + \boxed{\frac{k(k+1)}{2}} = \frac{(k+1)(k+2)}{2}$$

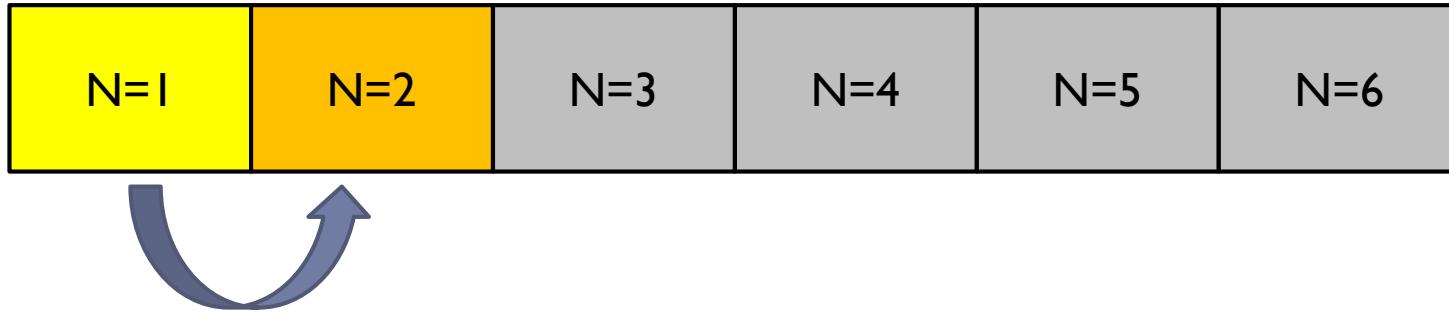
By assumption

Recursion & Induction

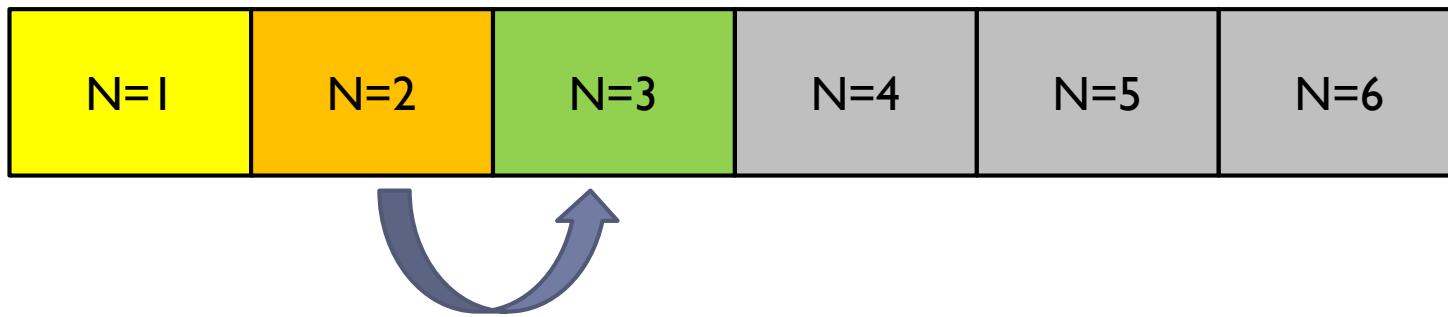
N=1	N=2	N=3	N=4	N=5	N=6
-----	-----	-----	-----	-----	-----



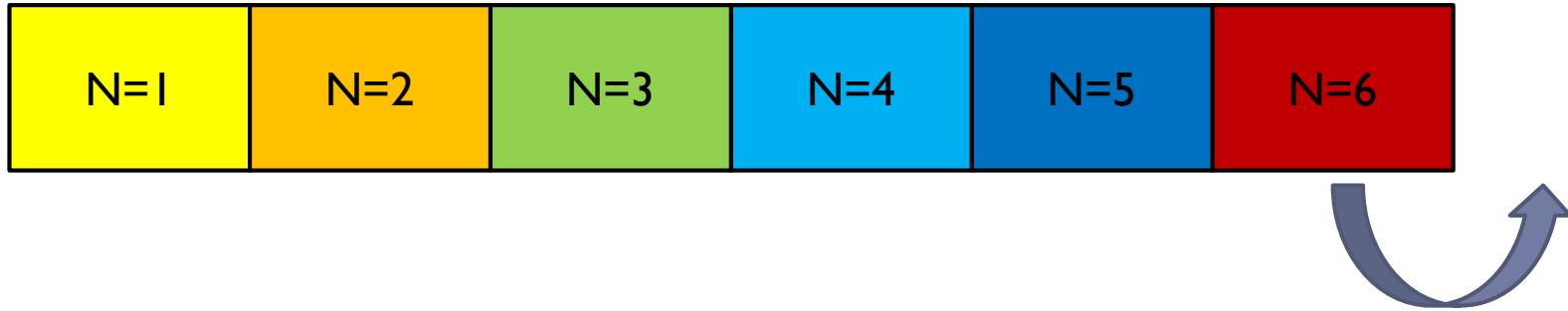
Recursion & Induction



Recursion & Induction



Recursion & Induction



Recursion & Induction

- ▶ 재귀는 수학적 귀납법과 매우 비슷하다
 - ▶ **But**, 증명 순서가 반대
 - ▶ $n = k-1$ 일 때 성립한다고 가정하면, $n = k$ 일 때 성립한다.
 - ▶ 당연히 초기조건이 있어야 함
- ▶ Factorial을 다시 생각해보면...



Recursion & Induction

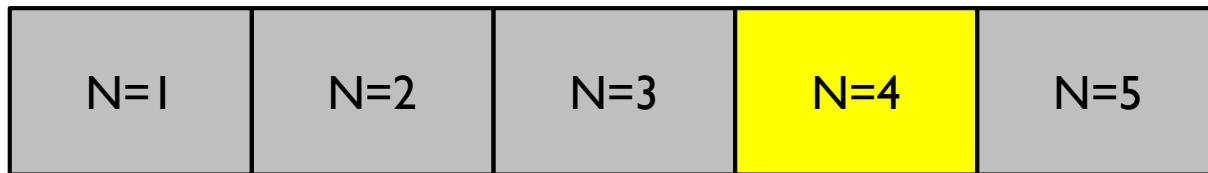
?

N=1	N=2	N=3	N=4	N=5
-----	-----	-----	-----	-----



Recursion & Induction

?



Recursion & Induction

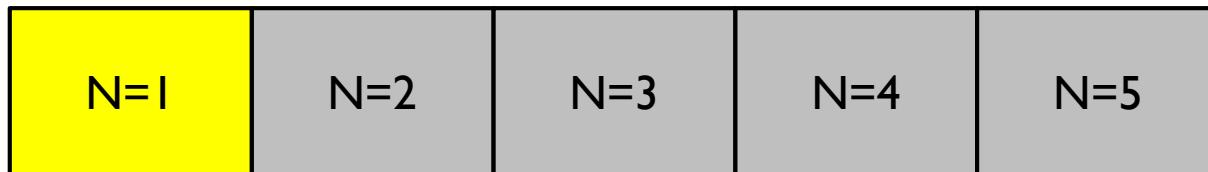
?

N=1	N=2	N=3	N=4	N=5
-----	-----	-----	-----	-----



Recursion & Induction

?



Recursion & Induction

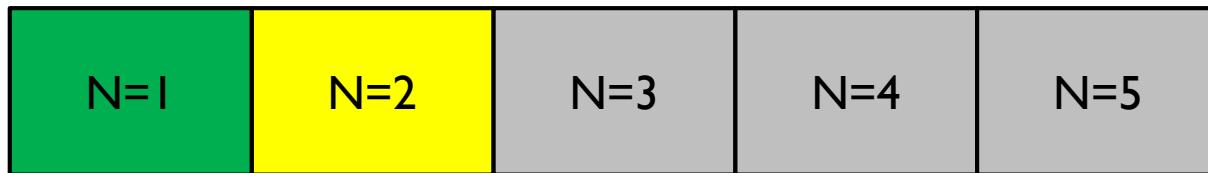
!!

N=1	N=2	N=3	N=4	N=5
-----	-----	-----	-----	-----



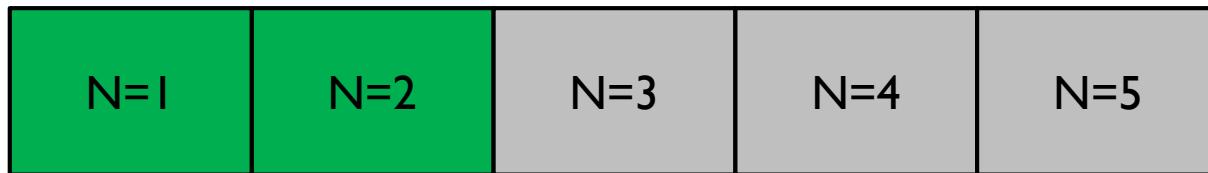
Recursion & Induction

?



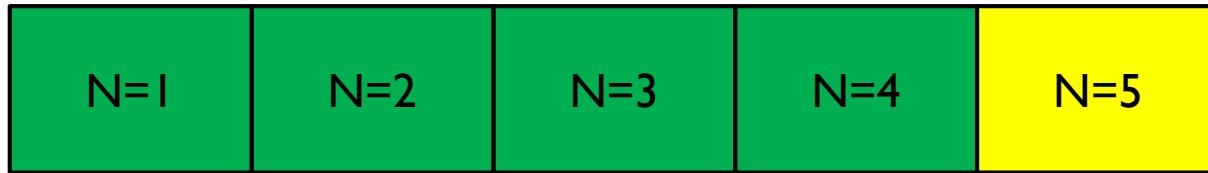
Recursion & Induction

!!



Recursion & Induction

?



Recursion & Induction

!!

N=1	N=2	N=3	N=4	N=5
-----	-----	-----	-----	-----



Recursion Implementation

- ▶ Function Decl {
 - Base Condition**
 - Process}

```
1: #include <csdio>
2: #include <algorithm>
3: using namespace std;
4: int get_fibo(int x){
5:     if(x == 0) return 1;
6:     else if(x == 1) return 1;
7:     else{
8:         int one = get_fibo(x-1);
9:         int two = get_fibo(x-2);
10:        return one + two;
11:    }
12: }
13: int main(){
14:     printf("%d",get_fibo(4));
15:     return 0;
16: }
```

Recursion Implementation

- ▶ 우선, 함수의 정의를 세운다
 - ▶ 귀납법에서 증명하고자 하는 명제를 세우는 것과 비슷
 - ▶ int get_fact(int x) : $x!$ 를 계산하여 return하는 함수
 - ▶ int get_fibo(int x) : F_x 를 계산하여 return하는 함수
- ▶ 그리고 귀납법처럼 생각하여 식을 세운다
 - ▶ $\text{return } x * \text{get_fact}(x-1)$
 - ▶ $\text{return get_fibo}(x-1) + \text{get_fibo}(x-2)$
- ▶ 참 쉽죠?



Recursion

- ▶ Recursion은 코딩을 함에 있어 첫 번째 난관
 - ▶ 극뽁! 해야 인생이 편해집니다
 - ▶ ~~착하게~~ ~~삽시다~~
- ▶ 여러가지 예제를 다루면서 익숙해 지는 것이 관건
- ▶ More Issue?



Recursion

- ▶ Divide & Conquer
- ▶ Dynamic Programming
- ▶ Back Tracking



Recursion

- ▶ **Divide & Conquer**
- ▶ Dynamic Programming
- ▶ Back Tracking



Divide & Conquer

- ▶ **Divide & Conquer (D&C)** is an important algorithm design paradigm based on multi-branched recursion. (wiki)
 - ▶ 말 그대로 분할해서 정ㅋ 벽ㅋ
 - ▶ 구체적으로 알아봅시다



Merge Sort

- ▶ Divide & Conquer의 가장 기본적인 형태
 - ▶ 일단 나누고, 각각에 대해서 작업을 수행한 다음
 - ▶ 하나로 합친다!
- ▶ 매우 중요함
 - ▶ Divide & Conquer의 대표적 모델로써 중요하고
 - ▶ $O(n \log n)$ 의 효율적인 정렬 알고리즘이라 중요하다
 - ▶ 응용이 될 가능성이 매우 큼



Merge Sort

▶ Merge Sort의 절차

- ▶ 1. 일단 두 부분으로 나눈다
- ▶ 2. 각각의 부분을 어떻게든 정렬한다.
- ▶ 3. 두 부분을 합친다



Merge Sort

1	4	3	2	4	7	1	2
---	---	---	---	---	---	---	---



Merge Sort

▶ I. 일단 나눈다

1	4	3	2	4	7	1	2
---	---	---	---	---	---	---	---



Merge Sort

▶ I. 일단 나눈다

1	4	3	2
---	---	---	---

4	7	1	2
---	---	---	---



Merge Sort

- ▶ 2. 어떻게든 둘을 정렬한다

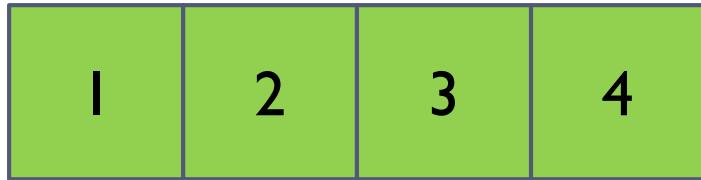
1	4	3	2
---	---	---	---

4	7	1	2
---	---	---	---



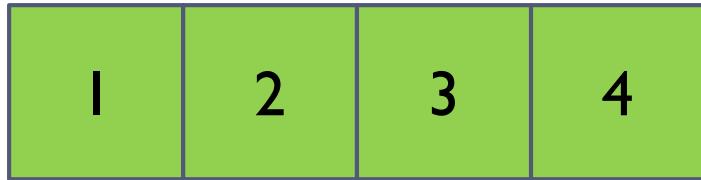
Merge Sort

- ▶ 2. 어떻게든 둘을 정렬한다



Merge Sort

- ▶ 3. 다시 하나로 합친다



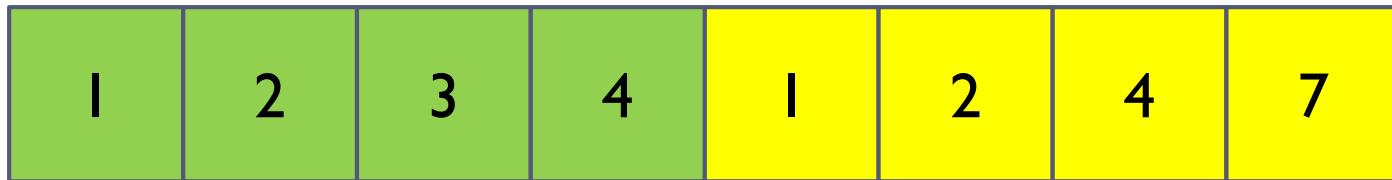
Merge Sort

- ▶ 3. 다시 하나로 합친다



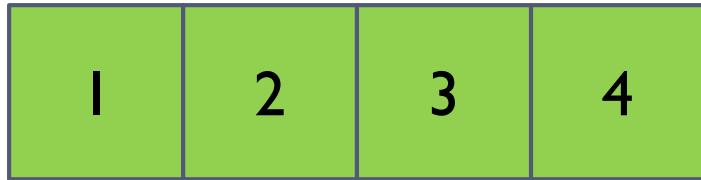
Merge Sort

- ▶ 3. 다시 하나로 합친다 ?????



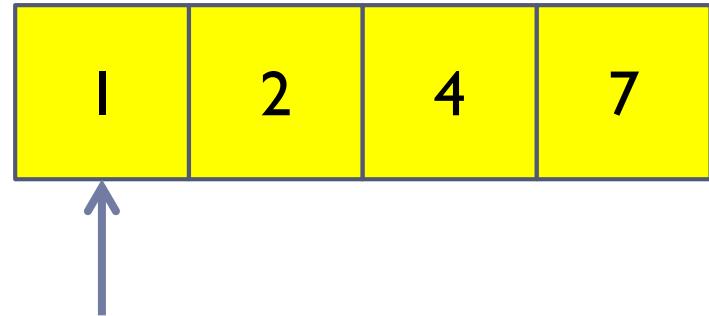
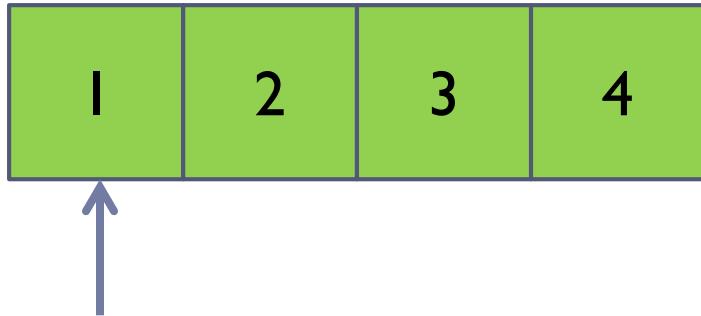
Merge Sort

- ▶ 3. 다시 하나로 합친다



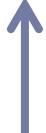
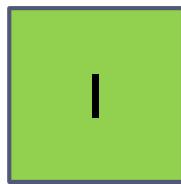
Merge Sort

- ▶ 3. 다시 하나로 합친다



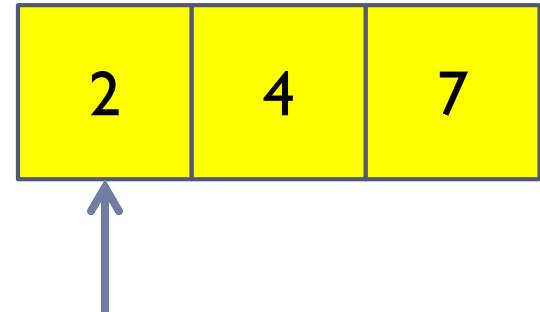
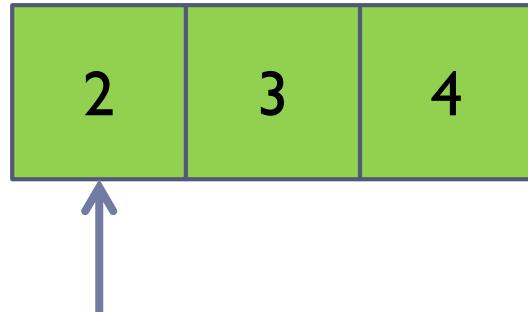
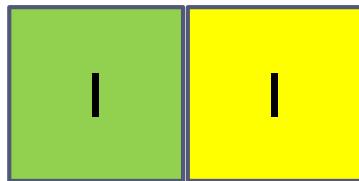
Merge Sort

- ▶ 3. 다시 하나로 합친다



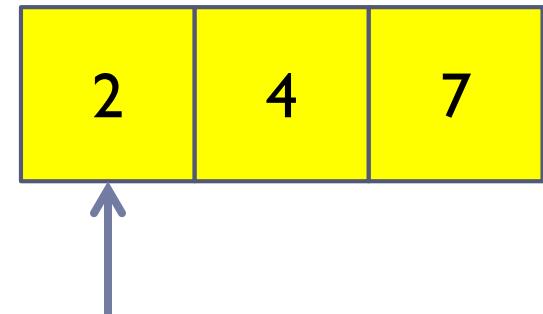
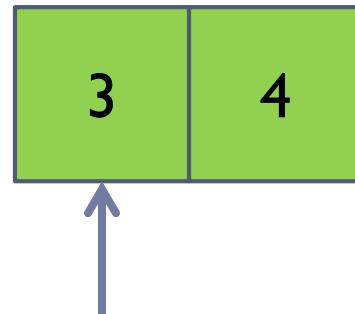
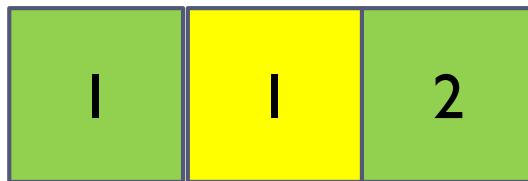
Merge Sort

- ▶ 3. 다시 하나로 합친다



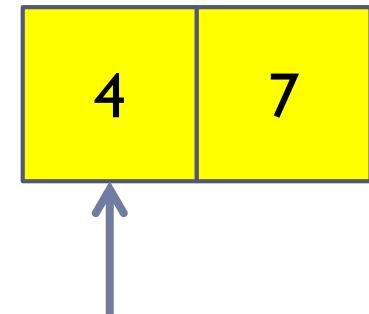
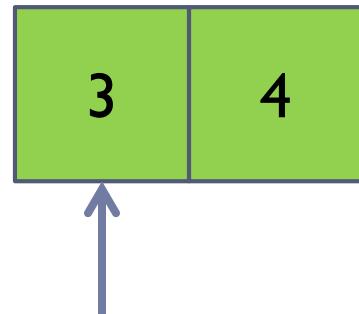
Merge Sort

- ▶ 3. 다시 하나로 합친다



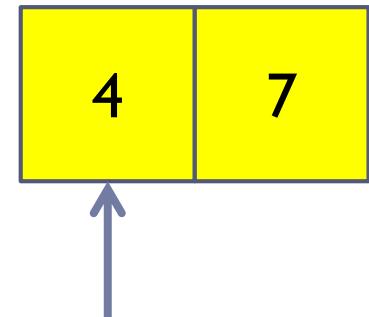
Merge Sort

- ▶ 3. 다시 하나로 합친다



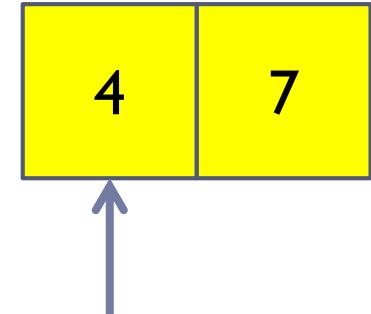
Merge Sort

- ▶ 3. 다시 하나로 합친다



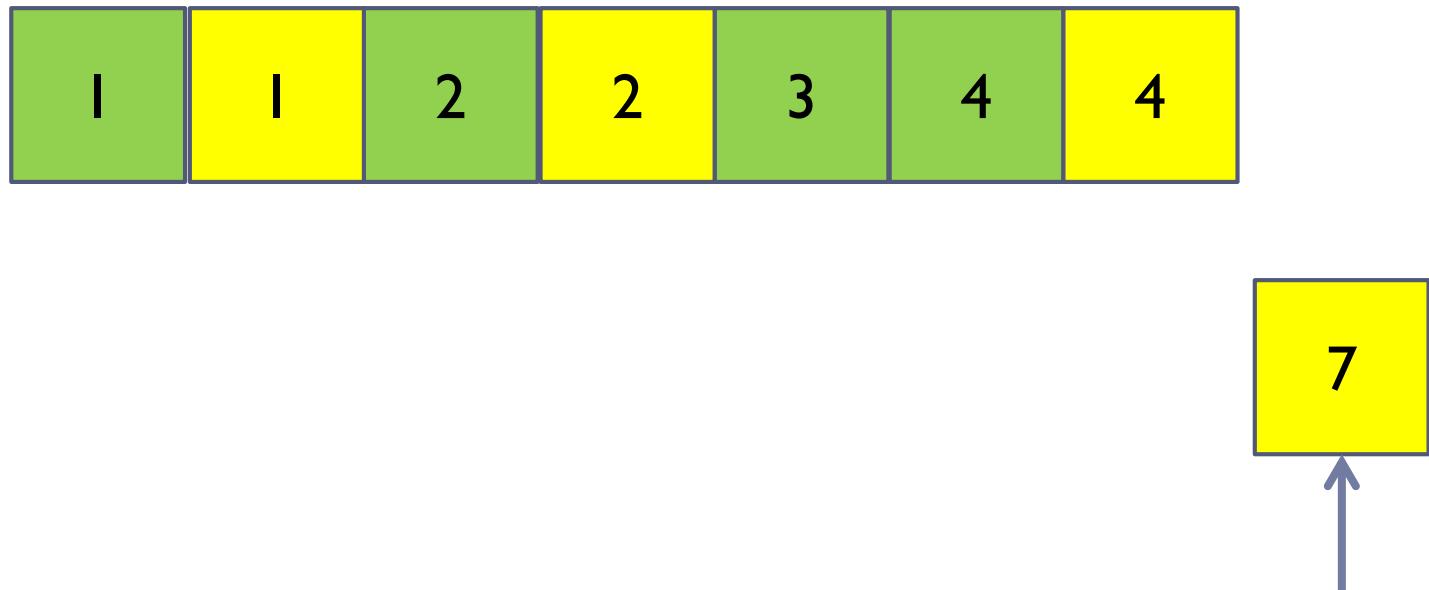
Merge Sort

- ▶ 3. 다시 하나로 합친다



Merge Sort

- ▶ 3. 다시 하나로 합친다



Merge Sort

- ▶ 3. 다시 하나로 합친다



Merge Sort

- ▶ 3. 다시 하나로 합친다 (!!)



Merge Sort

- ▶ 각각의 부분을 어떻게든 정렬한다??
→ 또 다시 Merge Sort로 정렬!

1	4	3	2	4	7	1	2
---	---	---	---	---	---	---	---



Merge Sort

- ▶ 각각의 부분을 어떻게든 정렬한다??
→ 또 다시 Merge Sort로 정렬!

1	4	3	2
---	---	---	---

4	7	1	2
---	---	---	---



Merge Sort

- ▶ 각각의 부분을 어떻게든 정렬한다??
→ 또 다시 Merge Sort로 정렬!

1	4
---	---

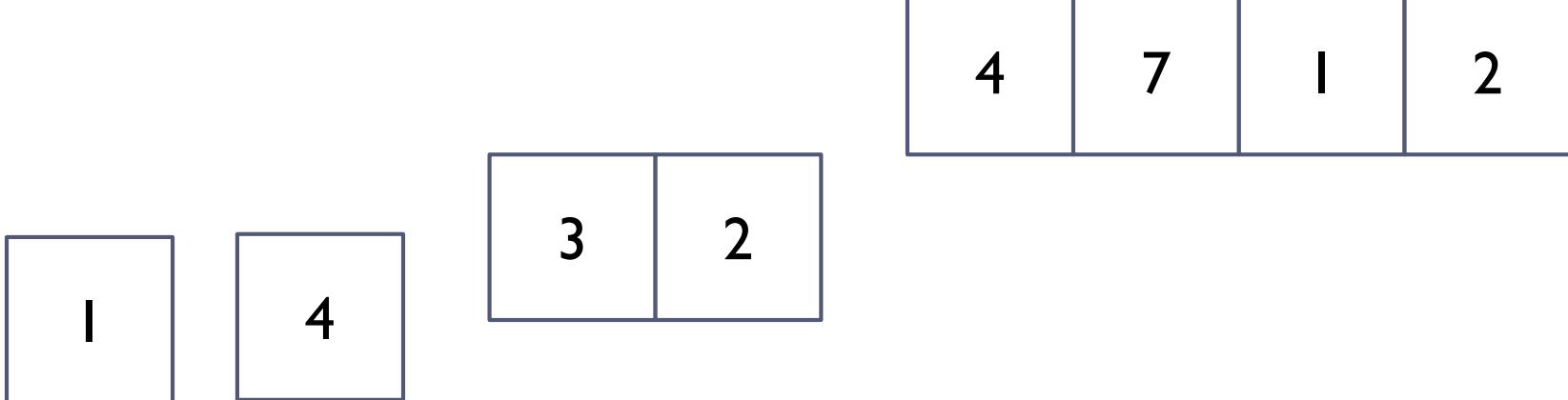
3	2
---	---

4	7	1	2
---	---	---	---



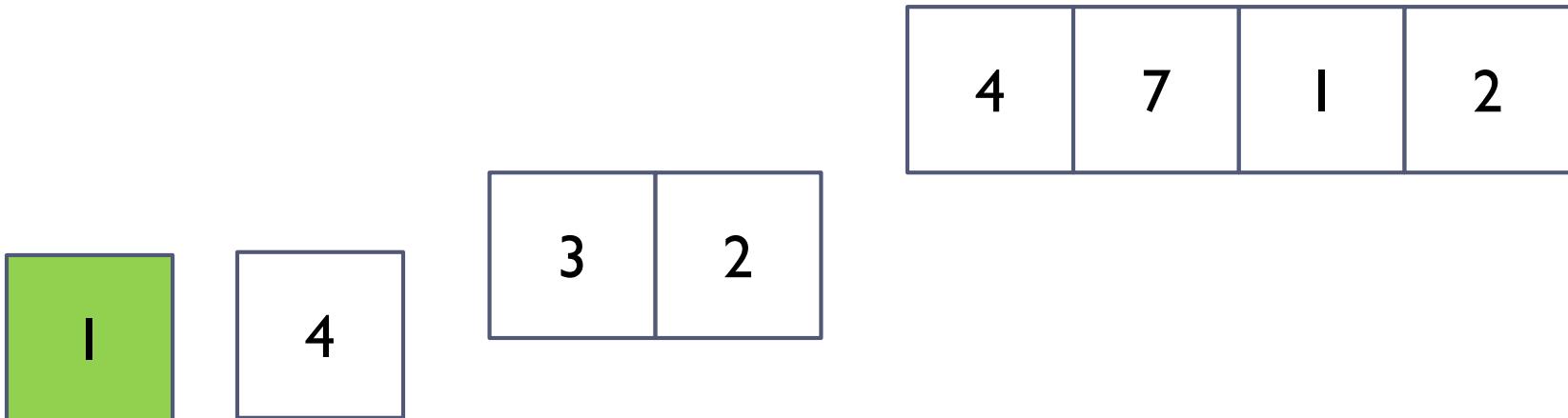
Merge Sort

- ▶ 각각의 부분을 어떻게든 정렬한다??
→ 또 다시 Merge Sort로 정렬!



Merge Sort

- ▶ 각각의 부분을 어떻게든 정렬한다??
→ 또 다시 Merge Sort로 정렬!



Merge Sort

- ▶ 각각의 부분을 어떻게든 정렬한다??
→ 또 다시 Merge Sort로 정렬!

1

4

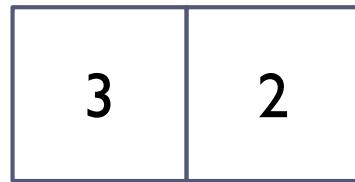
3	2
---	---

4	7	1	2
---	---	---	---



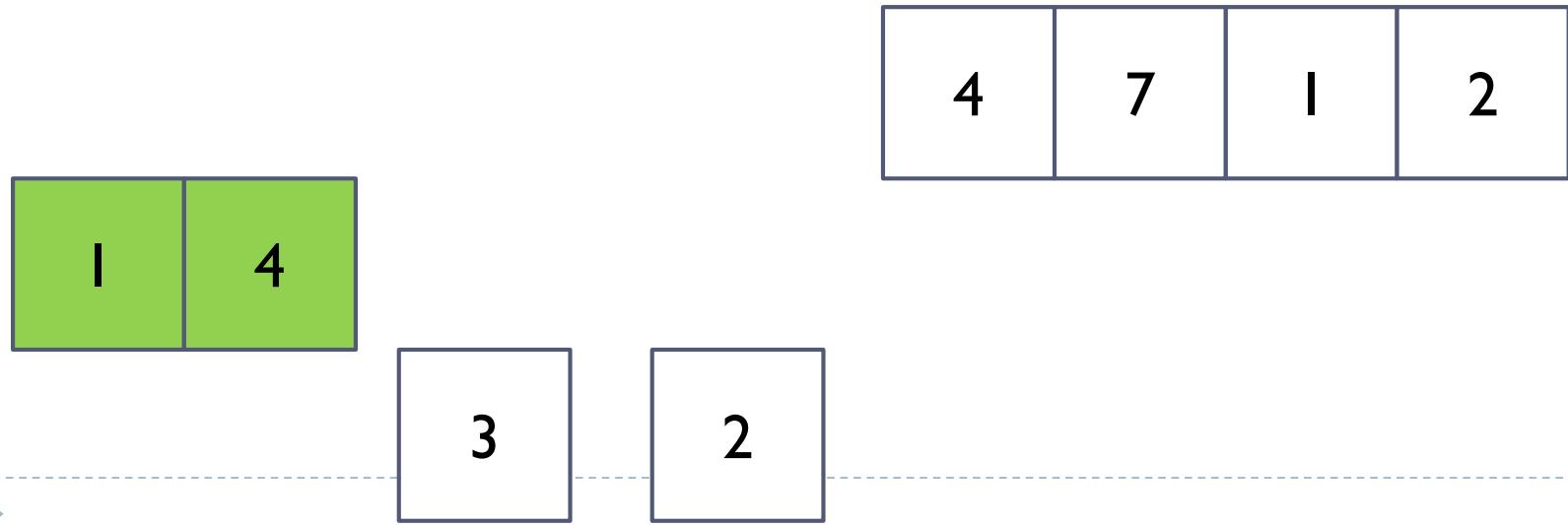
Merge Sort

- ▶ 각각의 부분을 어떻게든 정렬한다??
→ 또 다시 Merge Sort로 정렬!



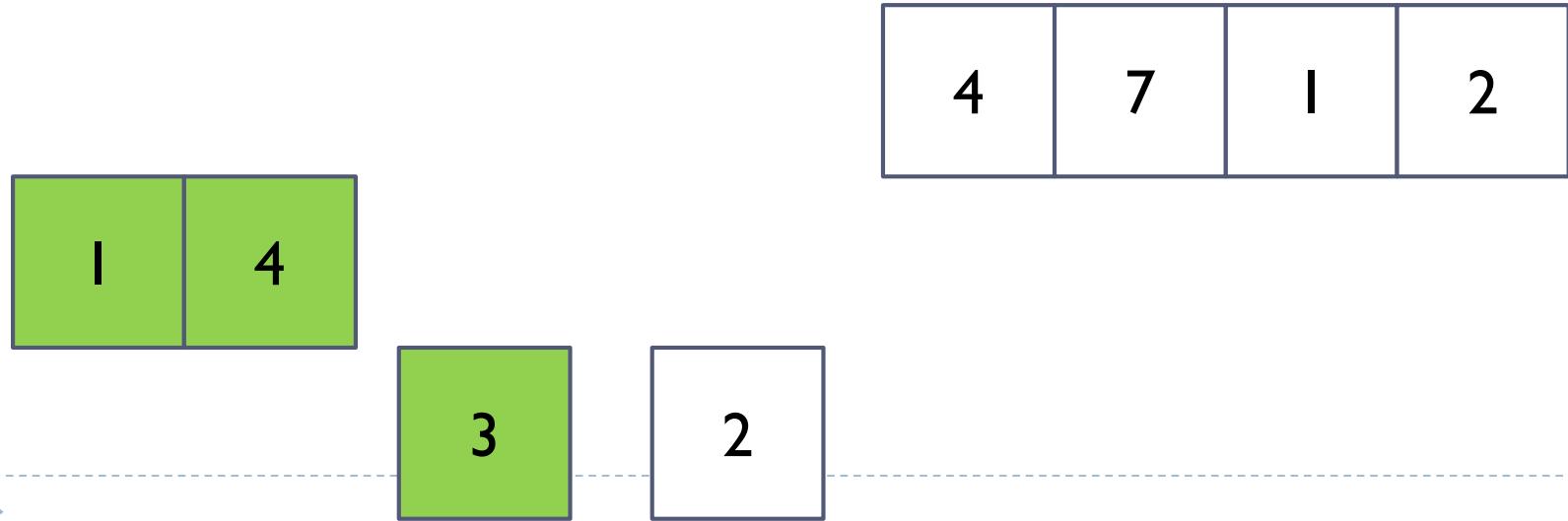
Merge Sort

- ▶ 각각의 부분을 어떻게든 정렬한다??
→ 또 다시 Merge Sort로 정렬!



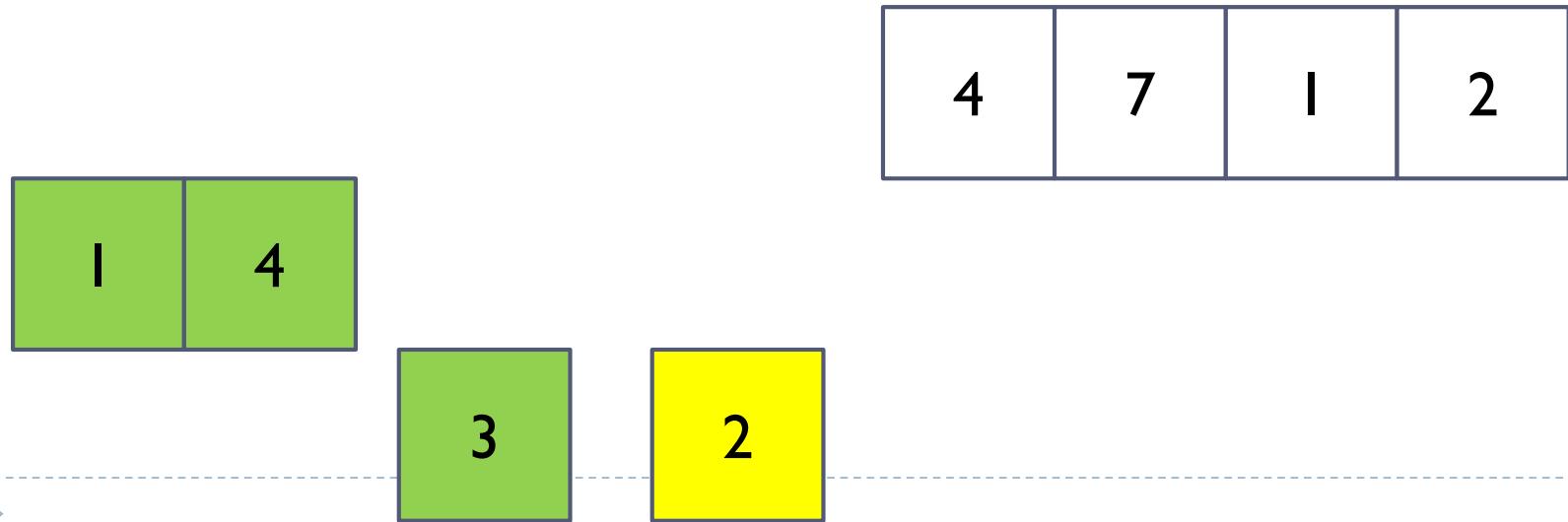
Merge Sort

- ▶ 각각의 부분을 어떻게든 정렬한다??
→ 또 다시 Merge Sort로 정렬!



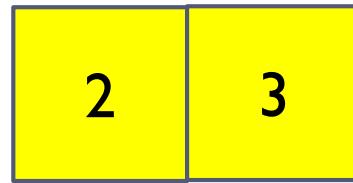
Merge Sort

- ▶ 각각의 부분을 어떻게든 정렬한다??
→ 또 다시 Merge Sort로 정렬!



Merge Sort

- ▶ 각각의 부분을 어떻게든 정렬한다??
→ 또 다시 Merge Sort로 정렬!



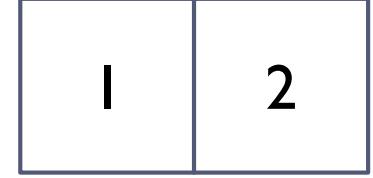
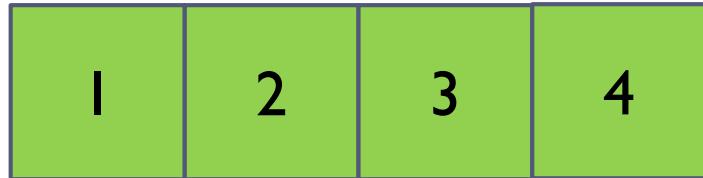
Merge Sort

- ▶ 각각의 부분을 어떻게든 정렬한다??
→ 또 다시 Merge Sort로 정렬!



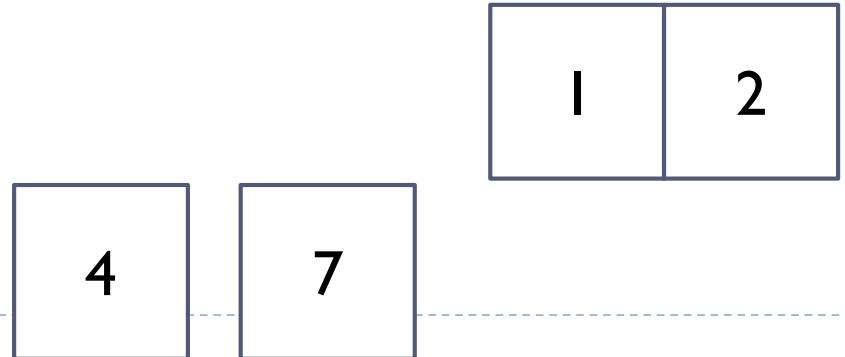
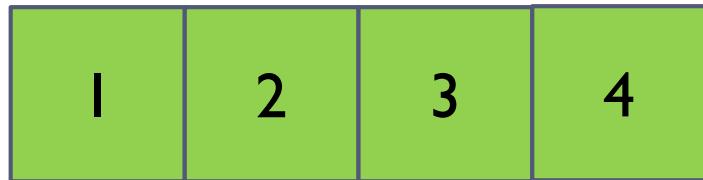
Merge Sort

- ▶ 각각의 부분을 어떻게든 정렬한다??
→ 또 다시 Merge Sort로 정렬!



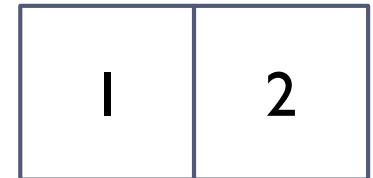
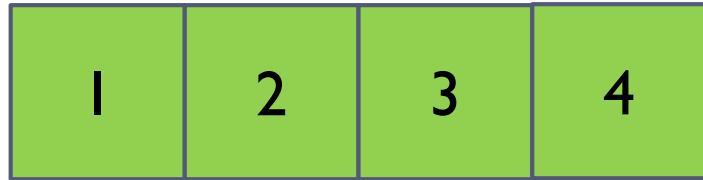
Merge Sort

- ▶ 각각의 부분을 어떻게든 정렬한다??
→ 또 다시 Merge Sort로 정렬!



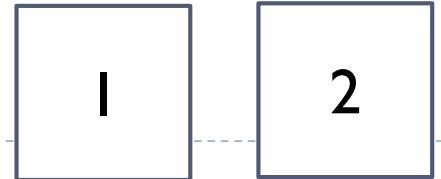
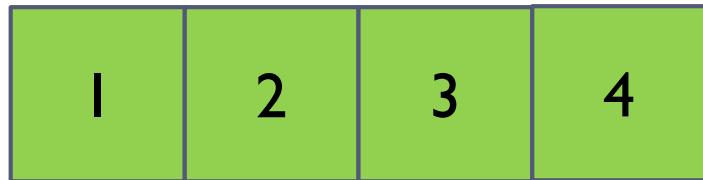
Merge Sort

- ▶ 각각의 부분을 어떻게든 정렬한다??
→ 또 다시 Merge Sort로 정렬!



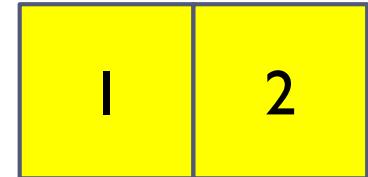
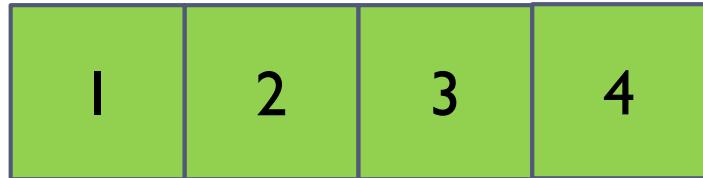
Merge Sort

- ▶ 각각의 부분을 어떻게든 정렬한다??
→ 또 다시 Merge Sort로 정렬!



Merge Sort

- ▶ 각각의 부분을 어떻게든 정렬한다??
→ 또 다시 Merge Sort로 정렬!



Merge Sort

- ▶ 각각의 부분을 어떻게든 정렬한다??
→ 또 다시 Merge Sort로 정렬!



Merge Sort

- ▶ 각각의 부분을 어떻게든 정렬한다??
→ 또 다시 Merge Sort로 정렬!



Divide & Conquer

```
4: int n, m;
5: int Data[30];
6: void MergeSort(int start, int end){
7:     if(start >= end) return;
8:     int mid = (start + end) / 2;
9:     MergeSort(start, mid);
10:    MergeSort(mid+1, end);
11:    // ----- merging -----
12:    int temp[30] = {0}, temp_len = 0; // temporary for merging
13:    int p1 = start, p2 = mid+1;
14:    while(p1 <= mid && p2 <= end){
15:        if(Data[p1] > Data[p2]){
16:            temp[temp_len++] = Data[p2];
17:            p2++;
18:        }
19:        else{
20:            temp[temp_len++] = Data[p1];
21:            p1++;
22:        }
23:    }
24:    if(p1 >= mid+1){ // right array has element
25:        while(p2 <= end){
26:            temp[temp_len++] = Data[p2];
27:            p2++;
28:        }
29:    }
30:    else if(p2 >= end+1){ // left array has element
31:        while(p1 <= mid){
32:            temp[temp_len++] = Data[p1];
33:            p1++;
34:        }
35:    }
36:    for(int i=start,j=0;i<=end;i++,j++) Data[i] = temp[j];
37:    return;
38: }
```



Divide & Conquer

```
39: int main(){
40:     int n;
41:     scanf("%d",&n);
42:     for(int i=1;i<=n;i++) scanf("%d",&Data[i]);
43:     MergeSort(1, n);
44:     for(int i=1;i<=n;i++) printf("%d ",Data[i]);
45:     printf("\n");
46:     return 0;
47: }
```



Divide & Conquer

- ▶ Merge Sort, Quick Sort가 대표적인 D&C.
 - ▶ Quick Sort 역시 마찬가지!



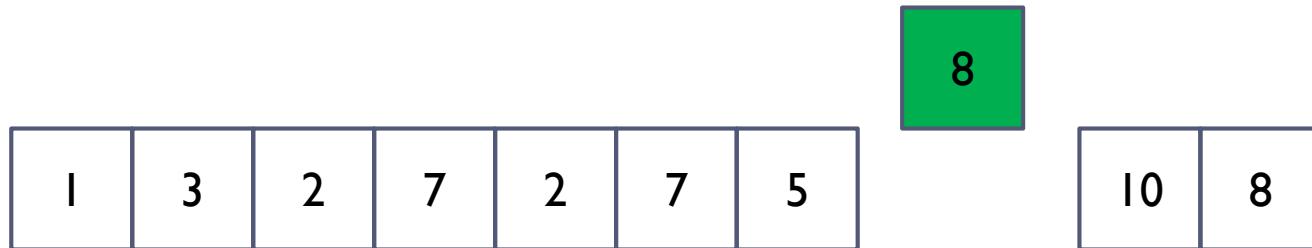
Divide & Conquer

- ▶ Merge Sort, Quick Sort가 대표적인 D&C.
 - ▶ Quick Sort 역시 마찬가지!



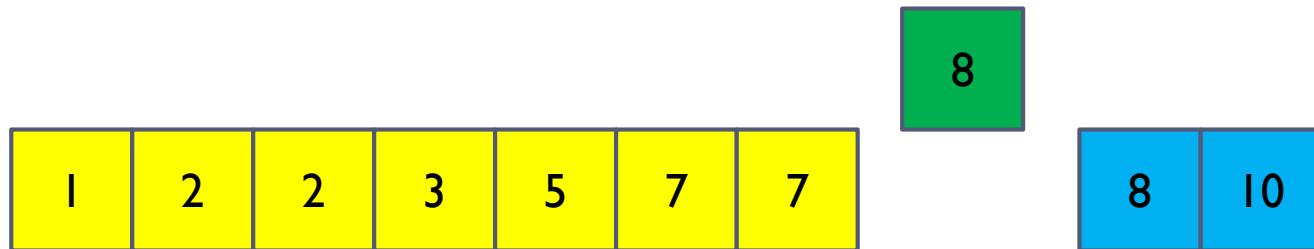
Divide & Conquer

- ▶ Merge Sort, Quick Sort가 대표적인 D&C.
 - ▶ Quick Sort 역시 마찬가지!



Divide & Conquer

- ▶ Merge Sort, Quick Sort가 대표적인 D&C.
 - ▶ Quick Sort 역시 마찬가지!



Divide & Conquer

- ▶ Merge Sort, Quick Sort가 대표적인 D&C.
 - ▶ Quick Sort 역시 마찬가지!



Divide & Conquer

- ▶ Merge Sort, Quick Sort가 대표
- ▶ Quick Sort 역시 마찬가지!



1	2	2	3	5	7	7	8	8	10
---	---	---	---	---	---	---	---	---	----



Divide & Conquer

- ▶ 문제를 나누어서 부분문제를 해결 후 전체문제 해결!
 - ▶ 부분 부분을 정렬한 후 전체 정렬
- ▶ D&C is **very** hard
 - ▶ D&C라고 판단하기가 어렵고
 - ▶ 판단을 하고 난 후에는 모델링이 난감합니다
 - ▶ For example,



Divide & Conquer

- ▶ 2차원에서 비교 불가능한 쌍의 개수를 찾아라
 - ▶ 비교가능? $(x_1, y_1) < (x_2, y_2)$ iff $x_1 < x_2 \& y_1 < y_2$
 - ▶ (1, 9) (7, 10) (4, 1) (3, 4) (5, 3) (6, 5) (2, 3) 라고 하자
 - ▶ 쉽게 하기 위해서 중복된 x좌표는 없다고 하자
- ▶ Indexed Tree를 사용하면 간단.
- ▶ Divide & Conquer로 해결!



Divide & Conquer

1	7	4	3	5	6	2
9	10	1	4	3	5	3



Divide & Conquer

1	2	3	4	5	6	7
9	3	4	1	3	5	10

☒순으로 정렬 (stable sort)



Divide & Conquer

1	2	3	4	5	6	7
9	3	4	1	3	5	10

이제 x를 제외하고 y만 가지면,



Divide & Conquer

9	3	4	1	3	5	10
---	---	---	---	---	---	----

이제 x를 제외하고 y만 가지면,



Divide & Conquer

9	3	4	1	3	5	10
---	---	---	---	---	---	----

자신보다 앞에 있는 숫자들보다 작으면 안됨!



Divide & Conquer

9	3	4	1	3	5	10
---	---	---	---	---	---	----

자신보다 앞에 있는 숫자들보다 작으면 안됨!



Divide & Conquer

9	3	4	1	3	5	10
---	---	---	---	---	---	----

이러한 쌍의 개수를 모두 찾자



Divide & Conquer

9	3	4	1	3	5	10
---	---	---	---	---	---	----

- I. 나누어서 각각의 쌍의 개수를 모두 찾는다
- II. 합치면서 새로 발생하는 쌍의 개수를 찾는다



Divide & Conquer

9	3	4	1	3	5	10
---	---	---	---	---	---	----

- I. 나누어서 각각의 쌍의 개수를 모두 찾는다
- II. 합치면서 새로 발생하는 쌍의 개수를 찾는다



Divide & Conquer

9	3	4
---	---	---

1	3	5	10
---	---	---	----

- I. 나누어서 각각의 쌍의 개수를 모두 찾는다
- II. 합치면서 새로 발생하는 쌍의 개수를 찾는다



Divide & Conquer

2개

9	3	4
---	---	---

0개

1	3	5	10
---	---	---	----

- I. 나누어서 각각의 쌍의 개수를 모두 찾는다
- II. 합치면서 새로 발생하는 쌍의 개수를 찾는다



Divide & Conquer

2개

0개

9	3	4	1	3	5	10
---	---	---	---	---	---	----

- I. 나누어서 각각의 쌍의 개수를 모두 찾는다
- II. 합치면서 새로 발생하는 쌍의 개수를 찾는다



Divide & Conquer

2개

6개

0개

9	3	4	1	3	5	10	

- I. 나누어서 각각의 쌍의 개수를 모두 찾는다
- II. 합치면서 새로 발생하는 쌍의 개수를 찾는다



Divide & Conquer

= 8개!

9	3	4	1	3	5	10
---	---	---	---	---	---	----

- I. 나누어서 각각의 쌍의 개수를 모두 찾는다
- II. 합치면서 새로 발생하는 쌍의 개수를 찾는다



Divide & Conquer

I. 은 단순히 나누어 재귀호출을 하면 됨

9	3	4	1	3	5	10
---	---	---	---	---	---	----

- I. 나누어서 각각의 쌍의 개수를 모두 찾는다
- II. 합치면서 새로 발생하는 쌍의 개수를 찾는다



Divide & Conquer

합치면서 새로 발생하는 것은 어떻게 찾나?

9	3	4	1	3	5	10
---	---	---	---	---	---	----

- I. 나누어서 각각의 쌍의 개수를 모두 찾는다
- II. 합치면서 새로 발생하는 쌍의 개수를 찾는다



Divide & Conquer

정렬이 되어있다면?

9	3	4	1	3	5	10
---	---	---	---	---	---	----

- I. 나누어서 각각의 쌍의 개수를 모두 찾는다
- II. 합치면서 새로 발생하는 쌍의 개수를 찾는다



Divide & Conquer

9	3	4	1	3	5	10
---	---	---	---	---	---	----



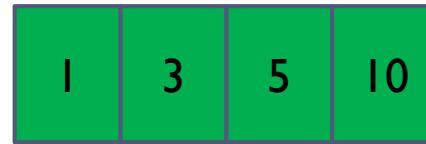
Divide & Conquer

9	3	4
---	---	---

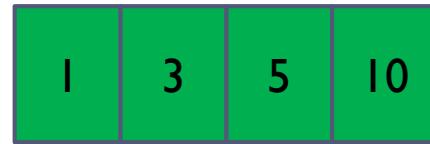
1	3	5	10
---	---	---	----



Divide & Conquer



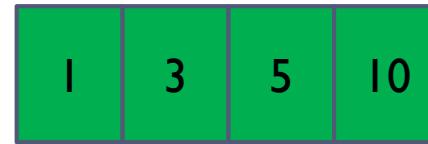
Divide & Conquer



왼쪽과 오른쪽이 정렬이 되어있다.
이제 Merging을 하면?



Divide & Conquer



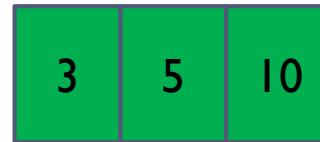
Divide & Conquer

3	4	9
---	---	---

3	5	10
---	---	----



Divide & Conquer



Divide & Conquer



Divide & Conquer

9

5 | 10

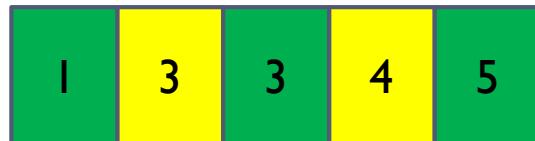
1 3 3 4



Divide & Conquer

9

10



Divide & Conquer

10



Divide & Conquer



Divide & Conquer



초록색은 오른쪽, 노란색은 왼쪽 묶음

초록색이 앞에 있으면 안됨!



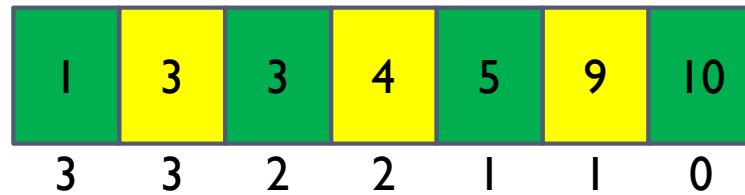
Divide & Conquer



초록색 입장에서 내 뒤에 노란색이 몇 개나 있나?



Divide & Conquer



초록색 입장에서 내 뒤에 노란색이 몇 개나 있나?



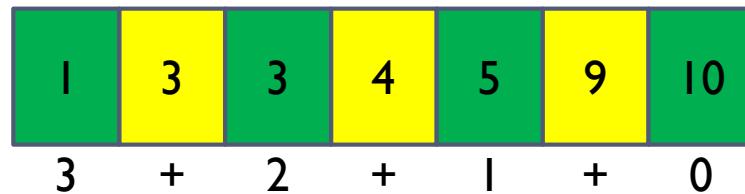
Divide & Conquer



초록색 입장에서 내 뒤에 노란색이 몇 개나 있나?



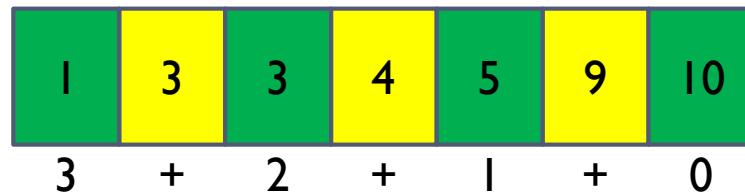
Divide & Conquer



초록색 입장에서 내 뒤에 노란색이 몇 개나 있나?



Divide & Conquer



합치면서 발생하는 쌍의 수가 6개!



Divide & Conquer

9	3	4	1	3	5	10
---	---	---	---	---	---	----

Overall procedure



Divide & Conquer

9	3	4	1	3	5	10
---	---	---	---	---	---	----



Divide & Conquer

9	3	4
---	---	---

1	3	5	10
---	---	---	----



Divide & Conquer

9

3 4

1 3 5 10



Divide & Conquer

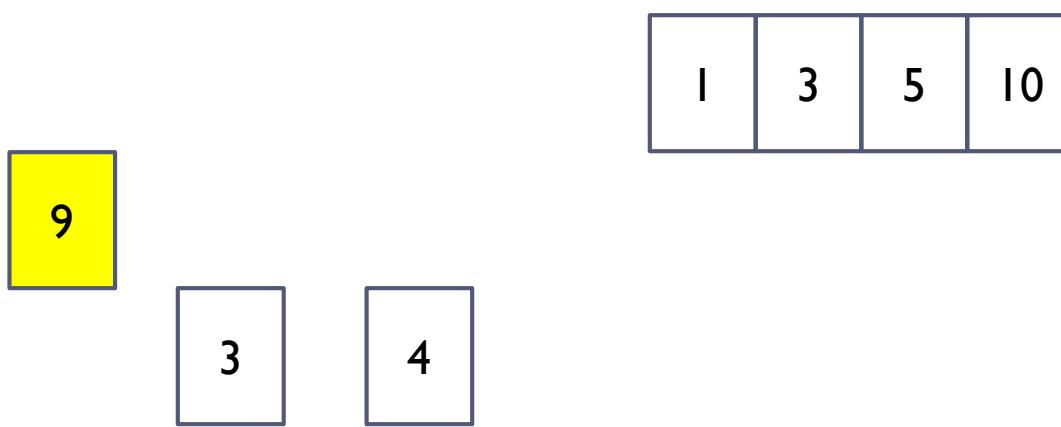
9

3	4
---	---

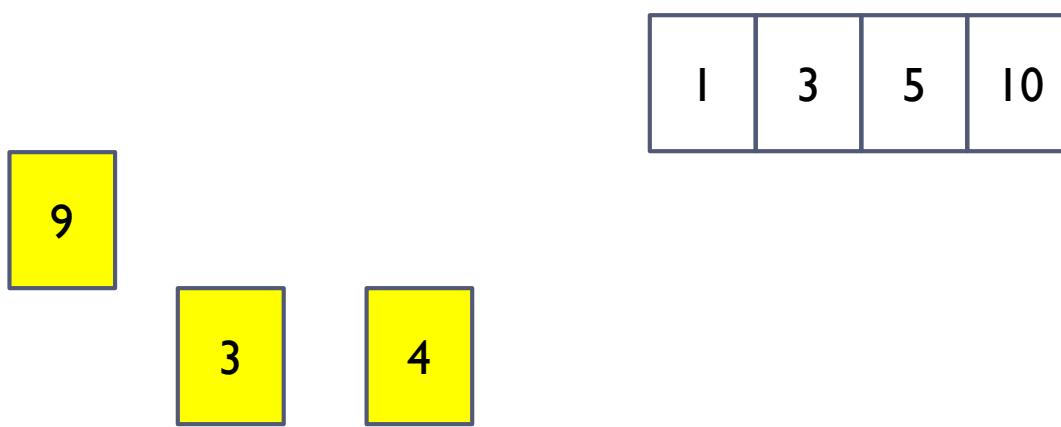
1	3	5	10
---	---	---	----



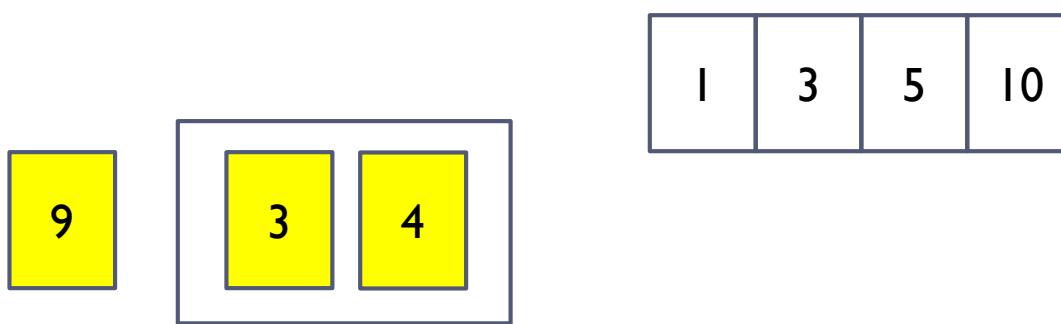
Divide & Conquer



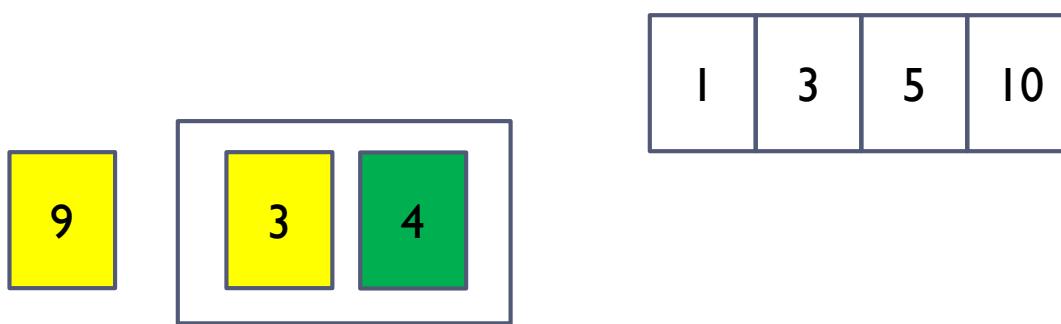
Divide & Conquer



Divide & Conquer



Divide & Conquer



Divide & Conquer

9

3 4

1	3	5	10
---	---	---	----

0



Divide & Conquer

9

3 4

1	3	5	10
---	---	---	----

0



Divide & Conquer

9

3 4

1 3 5 10

0



Divide & Conquer

9

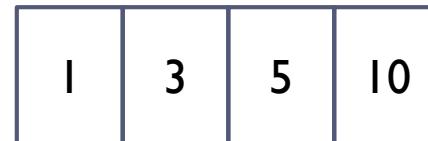
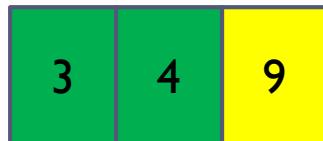
3 4

1 3 5 10

0



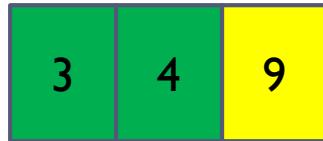
Divide & Conquer



0



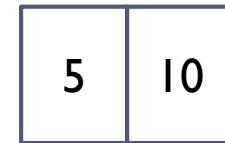
Divide & Conquer



$0 + 2$



Divide & Conquer



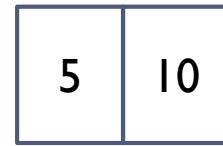
$0 + 2$



Divide & Conquer



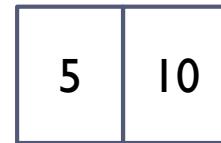
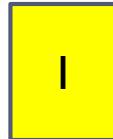
$0 + 2$



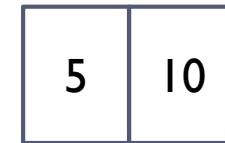
Divide & Conquer



$0 + 2$



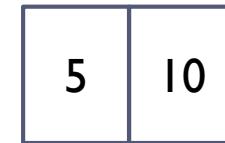
Divide & Conquer



$0 + 2$



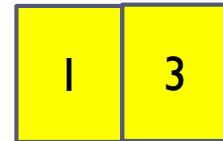
Divide & Conquer



$0 + 2 + 0$



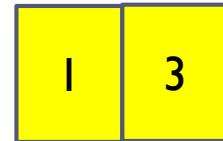
Divide & Conquer



$0 + 2 + 0$



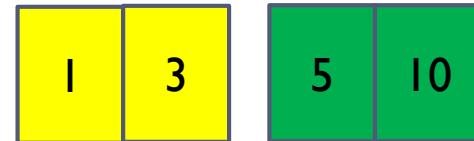
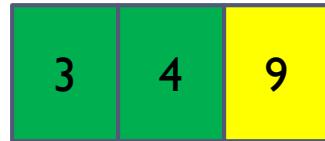
Divide & Conquer



0 + 2 + 0 + 0



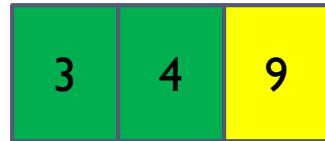
Divide & Conquer



0 + 2 + 0 + 0



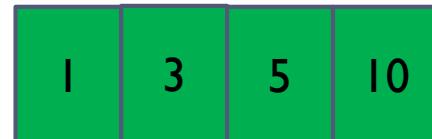
Divide & Conquer



$0 + 2 + 0 + 0 + 0$



Divide & Conquer



0 + 2 + 0 + 0 + 0



Divide & Conquer



0 + 2 + 0 + 0 + 0



Divide & Conquer



0 + 2 + 0 + 0 + 0



Divide & Conquer



$$0 + 2 + 0 + 0 + 0 + 6$$



Divide & Conquer

1	3	3	4	5	9	10
---	---	---	---	---	---	----

$$0 + 2 + 0 + 0 + 0 + 6 = 8 !$$



Divide & Conquer

- ▶ 3차원도 D&C로 풀릴 것입니다. (예전에 푼 기억이...)
- ▶ **D&C IS HARD!!**
 - ▶ 근데 그만큼 D&C로 어려운 문제 내기도 힘듭니다
 - ▶ 대표적인 D&C 문제들은 꼭 알아둡시다



Divide & Conquer

▶ Example

- ▶ N을 입력받아 | 2 3 ... (N-1) N (N-1) ... 3 2 | 출력
 - input = 4 → output = | 2 3 4 3 2 |
 - int Print (int s int n) : The output starts from s and peak point is n
- ▶ 1D array에서 최솟값 찾기
 - int Min(int array, int index) := The minimum value among array[0] ~ array[index]
- ▶ N개의 수를 입력 받아 연속부분 최대합을 출력
 - 함수 정의는 알아서 생각해보시길
- ▶ Hanoi's Tower
 - 판의 개수를 입력 받아 옮기는 과정을 출력한다
 - void Hanoi(int n, char a, char b, char c)
 - ▶ 첫 번째 기둥의 이름이 a, 두 번째 기둥의 이름이 b, 세 번째 기둥의 이름이 c이며, n개의 판이 a에 있을 때 옮기는 과정을 출력하는 함수

